Lemma. If $n \ge s \ge 0$ are integers and v(n) is the binary digit sum of n then $v(n-s) \ge v(n) - v(n\&s)$.

Proof. We assume n, s is the smallest counterexample (smallest n and for that n the smallest s) such that v(n-s) < v(n) - v(n & s).

If s = 0 then substituting this we get a contradiction:

$$v(n-s) < v(n) - v(n \& s)$$
$$v(n) < v(n) - v(0)$$
$$v(n) < v(n)$$

We may therefor assume that $s \geq 1$.

We will now assume that n, s have a binary digit in their binary expansion in common and denote this 2^b . We may define $n' = n - 2^b, s' = s - 2^b$. From this we get the following contradiction:

$$\begin{aligned} v(n-s) < &v(n) - v(n\&s) \\ v(n' + 2^b - s' - 2^b) < &v(n' + 2^b) - v(n'\&s' + 2^b) \\ v(n' - s') < &v(n') + 1 - v(n'\&s') - 1 \\ v(n' - s') < &v(n') - v(n'\&s') \end{aligned}$$

We may therefor assume that v(n&s) = 0 and so we must have n > s. Let us now assume that both n, s are even. We may define n = 2n' and s = 2s'. From this we get the following contradiction:

$$v(n-s) < v(n)$$

$$v(2n'-2s') < v(2n')$$

$$v(n'-s') < v(n')$$

If n is odd but s is even we may define n=2n'+1 and s=2s'. Since n>s we have 2n'+1>2s' leading to $n'\geq s'$. From this we get the following contradiction:

$$\begin{aligned} v(n-s) &< v(n) \\ v(2n'+1-2s') &< v(2n'+1) \\ v(n'-s') &+ 1 < v(n') + 1 \\ v(n'-s') &< v(n') \end{aligned}$$

If n is even and s is odd we may define n = n' + 1 and s = s' + 1. Since n > s we have n' > s'. If the lowest set bit in the binary representation of n is 2^l , $l \ge 1$, then we have v(n') = v(n) + l - 1. This is from the transition of 1 followed by

l 0's in n to a zero followed by l 1's in n'. n' and s' may have bits in common in bit positions 1...l-1. So we must have $0 \le v(n'\&s') \le l-1$ and hence $-l+1 \le -v(n'\&s') \le 0$. From this we get the following contradiction:

$$v(n-s) < v(n)$$

$$v(n'+1-s'-1) < v(n')-l+1$$

$$v(n'-s') < v(n')-v(n'\&s')$$